

Logical Characterizations of Local and Long-distance Phonological Agreement

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Characterizing phonological maps

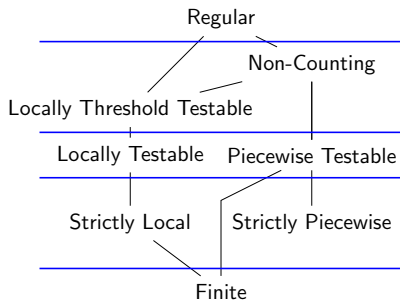
- **Goal:** identify the most restrictive computational complexity class that is sufficiently expressive for the set of possible phonological maps.
- Known computational complexity classes have multiple, converging characterizations.
 - finite state automata
 - formal language theory
 - algebraic
 - logical

Characterizing phonological maps

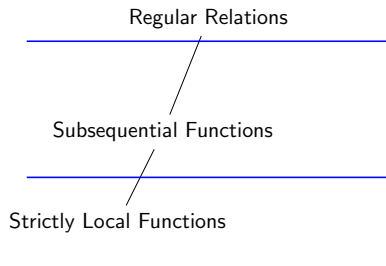
- Phonological maps are **regular relations** (Johnson 1972, Kaplan & Kay 1994)
- **Subregular** classes of functions appear to better characterize phonological maps (Mohri 1997; Chandlee et al. 2012; Gainor et al. 2012; Chandlee and Heinz 2012; Heinz and Lai 2013; Luo 2013; Payne 2013; Jardine, to appear)...
 - ...but not all classes are currently well understood.

Subregular hierarchies

Formal Languages



String-to-string maps



(Mohri 1997, Rogers & Pullum 2011, Rogers et al. 2013, Chandlee 2014)

Phonotactics

- Local phonotactics are **SL languages** (Heinz 2007, Heinz 2009)

(1) In a language with final devoicing:

- Allowable substrings: $\{ \times V, VV, V \times, \times D, \times T, DV, DD, TV, \dots \}$
- Prohibited substrings: $\{ D \times \}$

- Long-distance phonotactics are **SP** (Heinz 2010) or **TSL languages** (Heinz et al. 2011)

(2) In a language with sibilant harmony:

- Allowable subsequences: $\{ s \dots s, \int \dots \int, \dots \}$
- Prohibited subsequences: $\{ s \dots \int, \int \dots s \}$

Maps

- Local phonological maps are **SL functions** (Chandlee 2014, Chandlee et al. 2015)

$$(3) \quad D_{\times} \mapsto T_{\times}$$

- FST characterization
- Language-theoretic characterization
- **Hypothesis:** Long-distance maps are **SP functions**

$$(4) \quad s \dots f \mapsto s \dots s$$

Why use logic?

- Goal is to build on what is known of SL functions to identify the class of SP functions.
- Logic provides a natural extension from local to long-distance.

Examples of agreement

- Adjacent

(5) Korean (Lee and Pater 2008)
/p**a**pmul/ ↦ [p**a**mmul] 'rice water'

- Non-adjacent but bounded

(6) Ndonga (Viljoen 1973, Rose and Walker 2004)

a. /k**u**nila/ ↦ [k**u**nina] 'sow for'

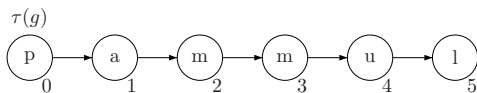
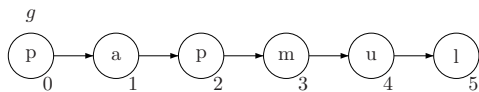
b. /n**i**kila/ ↦ [n**i**kila] 'season for'

- Unbounded

(7) Kikongo (Meinhof 1932, Odden 1994, Rose and Walker 2004)
/kud**u**mukisila/ ↦ [kud**u**mukisina] 'to cause to jump for'

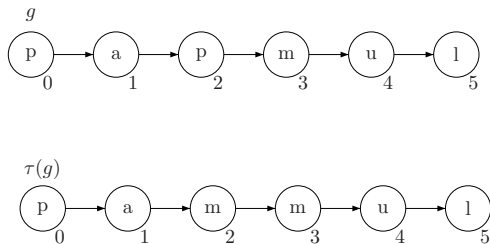
String graphs

$\text{papmul} \mapsto \text{pammul}$



String graphs

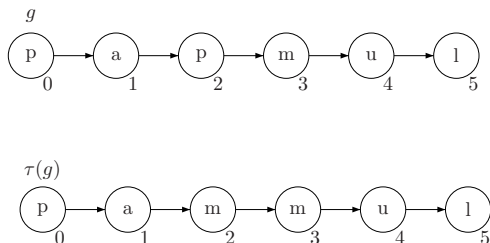
papmul \mapsto pammul



- Set of nodes: $\{0, 1, 2, 3, 4, 5\}$
- Successor relation: $S(0) = 1, S(1) = 2$, etc.
(also written $1 \triangleleft 2$)

String graphs

papmul \mapsto **pammul**



- Labeling function: $\ell(0) = p$, $\ell(1) = a$, etc.

Natural class predicates

- $V(x) = \text{TRUE}$ iff $\ell(x) = \mathbf{a} \vee \ell(x) = \mathbf{u} \vee \dots$
- $N(x) = \text{TRUE}$ iff $\ell(x) = \mathbf{n} \vee \ell(x) = \mathbf{m} \vee \dots$
- $C(x) = \text{TRUE}$ iff $\neg N(x) \wedge \neg V(x)$

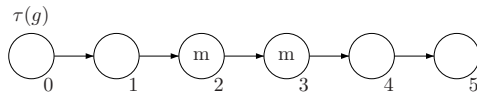
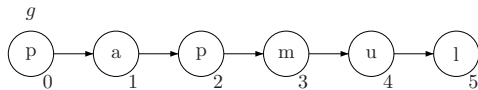
Substring predicates

- $CN(x) = (\exists y) [C(x) \wedge N(y) \wedge x \triangleleft y]$

Graph transduction

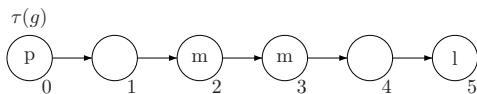
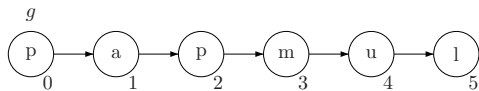
- Predicates define the nodes and labels of the output graph in terms of the input graph (Engelfriet and Hooageboom 2001)

$$\varphi_N^0(x) = N(x) \vee CN(x)$$



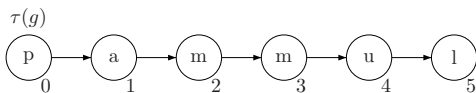
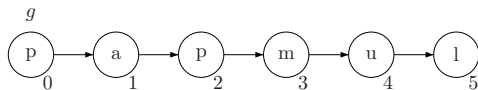
Graph transduction

$$\varphi_C^0(x) = C(x) \wedge \neg CN(x)$$



Graph transduction

$$\varphi_V^0(x) = V(x)$$



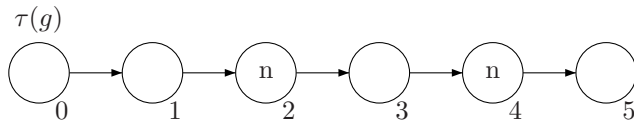
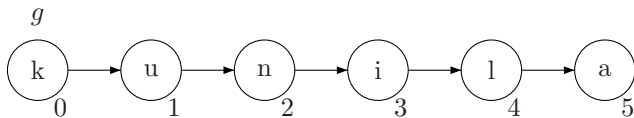
Non-adjacent but bounded

kunila \mapsto **kunina**

$$\text{NVL}(x) = (\exists y, z) [N(y) \wedge V(z) \wedge L(x) \wedge y \triangleleft z \wedge z \triangleleft x]$$

Non-adjacent but bounded

$$\varphi_N^0(x) = N(x) \vee NVL(x)$$

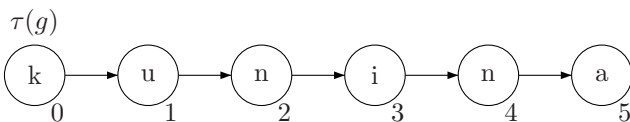
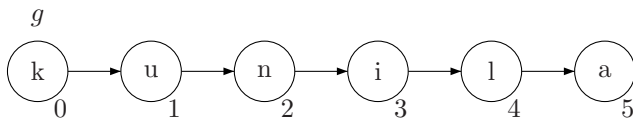


Non-adjacent but bounded

$$\varphi_L^0(x) = L(x) \wedge \neg \text{NVL}(x)$$

$$\varphi_V^0(x) = V(x)$$

$$\varphi_D^0(x) = D(x)$$



Unbounded

tukunidi \mapsto tukunini

$$\text{NVL}(x) = (\exists y, z) [\text{N}(y) \wedge \text{V}(z) \wedge \text{L}(x) \wedge y \triangleleft z \wedge z \triangleleft x]$$

Unbounded

tunikidi \mapsto **tunikini**

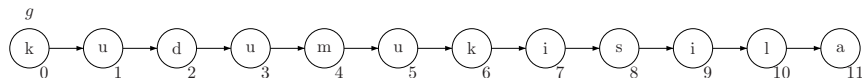
$$\begin{aligned} N \dashv\vdash L(x) &= (\exists v, w, y, z) [N(y) \wedge L(x) \\ &\quad \wedge y \triangleleft z \wedge z \triangleleft w \wedge w \triangleleft v \wedge v \triangleleft x] \end{aligned}$$

Unbounded

kudumukisila \mapsto kudumukisina

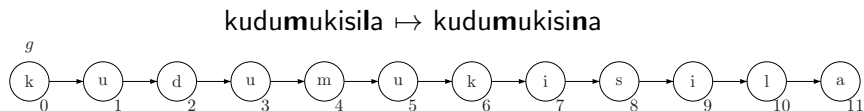
$$N \text{---} \text{---} \text{---} \text{---} L(x) = (\exists s, t, u, v, w, y, z) [N(y) \wedge L(x) \\ \wedge y \triangleleft z \wedge z \triangleleft w \wedge w \triangleleft v \wedge v \triangleleft u \wedge u \triangleleft t \wedge t \triangleleft s \wedge s \triangleleft x]$$

Instead of successor, **precedence**



- Successor: $x \triangleleft y$
 - $S(0) = 1, S(1) = 2, \dots$
- Precedence: $x < y$
 - $P(1) = \{0\}, P(2) = \{0, 1\}, \dots, P(10) = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$

Instead of substring, **subsequence**



$$NL(x) = (\exists y) [N(y) \wedge L(x) \wedge y < x]$$

Categories

Local	Long-distance
Adjacent	Non-adjacent
Bounded	Unbounded
SL	SP

More examples of bounded LDA

- (8) **Transvocalic:** Anywa, Basaa, Bemba, Bukusu, Dime, Gaagudju, Herero, Ila, Kipare, Koyra, Kwanyama, Lamba, Luba, Mayali, Mwiini, Nyangumarta, Pangwa, Pare, Pende, Punu, Ruund, Sawai, Shambaa, Sudanese, Suku, Tonga, Yabem, Zayse, ...
- (9) **Co-occurrence restrictions within roots or morphemes:** Adhola, Alur, Anywa, Basque, Chaha, Chol, Ganda, Hausa, Ineseno Chumash, Izere, Kalasha, Karaim, Komi-Permyak, Kukuya, Luo, Malto, Ndebele, Ngizim, Pengo, Pohnpeian, Shilluk, Siwi, Teke-Gabon, Tiene, Yucatec Mayan, Zulu, ...

(see Shaw (1991), Odden (1994), Hansson (2001), Rose & Walker (2004), and references therein)

Domain restrictions

(10) Malto (Mahapatra 1979, Hansson 2001): $\{t, d\}$ and $\{t, d\}$ cannot co-occur in tautomorphic CVC

- SL function that targets substrings of length 3?

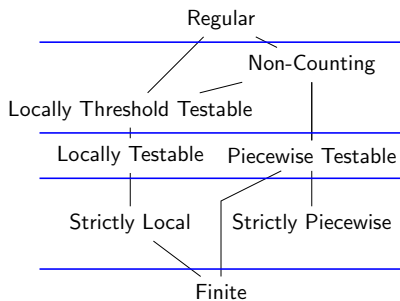
$$\begin{array}{ll} ta\downarrow \mapsto tat & taat\downarrow \mapsto taat \\ da\downarrow \mapsto dad & dana\downarrow \mapsto danad \end{array}$$

- Or SP function that targets subsequences of length 2 that will never 'see' inputs longer than 3?

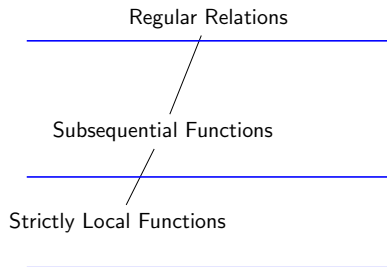
$$\begin{array}{ll} ta\downarrow \mapsto tat & taat \mapsto taat \\ da\downarrow \mapsto dad & danad \mapsto danad \end{array}$$

Subregular hierarchies

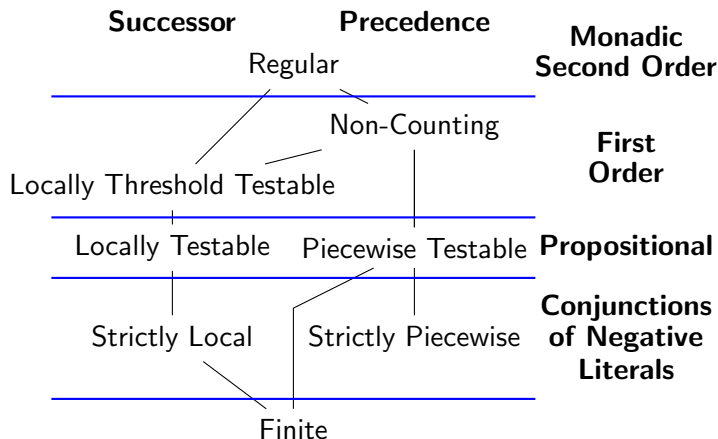
Formal Languages



String-to-string maps

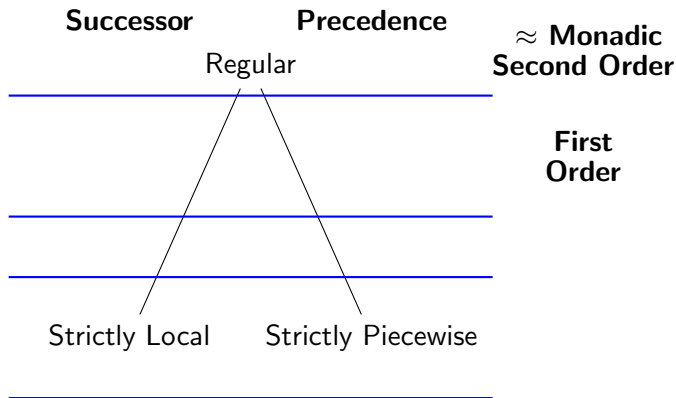


Subregular hierarchy of languages



(Rogers & Pullum 2011, Rogers et al. 2013)

Subregular hierarchy of functions



First Order Logic

- \forall, \exists
- \wedge, \vee, \neg
- $\rightarrow, \leftrightarrow$
- x, y, z, \dots
- $P(x), R(x, y), \dots$

'Sub'-First Order Logic

Proposal: SL functions are quantifier-free, unary, monotone graph transductions (Chandlee & Lindell, in prep)

- ~~\forall, \exists~~
- \wedge, \vee, \neg
- $\rightarrow, \leftrightarrow$
- x, y, z, \dots
- $P(x), R(x, y), \dots$

Next steps

- Extend to dissimilation maps
- Logical and FST characterizations of SP functions
- Interactions of SL and SP maps?

Conclusions

- Characterizing computational complexity classes from a variety of perspectives leads to a more complete understanding of what the class is.
- Identifying the right class (or combination of classes) for phonological maps likewise leads to insights into what phonology is (and isn't).
- From this perspective, the distinction between local and long-distance phonology is one of representation, not computational complexity.

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SP FSTs?

(11) Benchnon sibilant harmony (Hayward 1988, Hansson 2001)

- a. /s^jap-s/ \mapsto [s^japs] 'make wet'
 b. /ʃir-s/ \mapsto [ʃirʃ] 'bring near'

V:V, s:s, C:C C:C, V:V, ʃ:ʃ, s:ʃ

